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On the Ziv-Merhav estimator

Quantissima in the Serenissima V
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Project: *ENIC - Entropy, information and control of complex processes*

Joint work with:

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Entropies

(Shannon) entropy rate means:

$$h(\mathbb{P}) = \lim_{n \rightarrow \infty} -\frac{1}{n} \sum_{a \in \mathcal{A}^n} \mathbb{P}[a] \log \mathbb{P}[a]$$

Cross entropy rate – when existing – means:

$$h^c(\mathbb{P}|\mathbb{Q}) = \lim_{n \rightarrow \infty} -\frac{1}{n} \sum_{a \in \mathcal{A}^n} \mathbb{P}[a] \log \mathbb{Q}[a]$$

Relative entropy rate (or Kullback-Leibler divergence):

$$h^f(\mathbb{P}|\mathbb{Q}) = h^c(\mathbb{P}|\mathbb{Q}) - h(\mathbb{P})$$

Universal estimators: no complete knowledge of the measures!

We will always assume that \mathbb{P} and \mathbb{Q} are shift-invariant with respect to the left-shift operator $Tx_k := x_{k+1}$ with $x \in \mathcal{A}^{\mathbb{N}}$.

The waiting time problem

- **Kac's Lemma (1947):** given $R_n(x) := \inf\{\ell \geq 1 \mid x_{\ell+1}^{\ell+n} = x_1^n\}$, $\forall a \in \mathcal{A}^n$ we have

$$\log \mathbb{E}_a[R_n] := \log \int R_n(x) d\mathbb{P}(x \mid x_1^n = a) = -\log \mathbb{P}[a];$$

- **Wyner-Ziv (1989) and Ornstein-Weiss (1993):** if \mathbb{P} is *ergodic*, then for \mathbb{P} -almost every x we have

$$\lim_{n \rightarrow \infty} \frac{\log R_n(x)}{n} = h(\mathbb{P}).$$

*No strong assumptions on the measure:
access to the entropy just looking at the data!*

Similarly, $W_n(x, y) := \inf_{\ell \geq 1} \{y_\ell^{\ell+n-1} = x_1^n\}$ for two sequences $x \sim \mathbb{P}, y \sim \mathbb{Q}$ and we expect

$$W_n(x, y) \sim e^{h^c(\mathbb{P}|\mathbb{Q}) \cdot n}, \quad \text{where} \quad h^c(\mathbb{P}|\mathbb{Q}) = h(\mathbb{P}) + h^f(\mathbb{P}|\mathbb{Q})$$

for “sufficiently well-behaved” measures:

- **Kontoyiannis (1998)**: \mathbb{Q} is ψ -mixing or summably ϕ -mixing;
- **Cristadoro, Degli Esposti, Jakšić and Raquépas (2023)**: decoupling perspective for \mathbb{Q} and ergodicity for \mathbb{P} .

Computational drawback: given x, y *finite* sequences, even for small values of n , you may scroll all the data concerning x without having seen the prefix $y_1^n \dots$

⇒ **Dual quantity**: longest-match length! [Kontoyiannis (1998)]

$$\Lambda_n(x, y) := \sup_{\ell \geq 1} \{W_\ell(x, y) \leq n - \ell + 1\} \\ \leq n$$

E.g. $x_1^{10} = \underbrace{101\ 1100100}$
 $y_1^\infty = 01 \underbrace{101}_{x_1^3} \underbrace{00010}_{x_1^5} | 11120100\dots$

$$\lim_{n \rightarrow \infty} \frac{\log W_n(x, y)}{n} = \lim_{n \rightarrow \infty} \frac{\log n}{\Lambda_n(x, y)} = h^c(\mathbb{P}|\mathbb{Q}) \quad \text{for } (\mathbb{P} \otimes \mathbb{Q})\text{-almost every } (x, y)$$

...BUT

Ziv-Merhav: a waste-free estimator!

Heuristics: Iteration!

Compute $\Lambda_n(x, y)$ at different (**dependent**) starting locations throughout x_1^n and take the average!

$x = 01100101000102011101001000210 \dots,$
 $y = 010\underbrace{001011}_{x_4^8}1010 \underbrace{011}_{x_1^3} 1001000122021 \dots,$

$$x_1^{24} = 011|00101|00010|2|011101001|0 \implies c_{24}^{ZM}(x|y) = 6$$

$c_n^{ZM}(x|y)$ is the number of words in the sequential parsing of x_1^n using longest-matches with y_1^n .

We expect: $\lim_{n \rightarrow \infty} Q_n^{ZM}(x, y) := \lim_{n \rightarrow \infty} \frac{c_n^{ZM}(x|y) \log n}{n} = h^c(\mathbb{P}|\mathbb{Q}) \quad (\mathbb{P} \otimes \mathbb{Q})\text{-a.s.}$

Why? Define $\ell_1 := \Lambda_n(x, y)$ and $\ell_j := \Lambda_n(T^{\ell_1 + \dots + \ell_{j-1}}x, y)$ with $\sum_{j=1}^{c_n^{ZM}} \ell_j = n$, then

$$\frac{c_n^{ZM} \log n}{n} = \frac{\log n}{\frac{1}{c_n^{ZM}} \sum_{j=0}^{c_n^{ZM}-1} \Lambda_n(T^{\ell_1 + \dots + \ell_{j-1}}x, y)} \stackrel{?}{\approx} \frac{\log n}{\Lambda_n(x, y)}$$

mZM estimator and decoupling perspective

Step I: do the best in the spirit of the original proof

Theorem (J. Ziv and N. Merhav, 1993)

Let \mathbb{P}, \mathbb{Q} be two ergodic multi-level Markov measures. Then

$$\lim_{n \rightarrow \infty} Q_n^{\text{ZM}}(x, y) = h^c(\mathbb{P}|\mathbb{Q})$$

for $(\mathbb{P} \otimes \mathbb{Q})$ -almost every $(x, y) \in \mathcal{A}^{\mathbb{N}} \times \mathcal{A}^{\mathbb{N}}$.

*Which properties of irreducible multi-level Markov chains
do they really use?*

“Strong” decoupling (for the auxiliary parsing)

[no gaps: $\tau = 0$ in **SLD**]

There exists $(k_n)_{n=1}^{\infty} \subset [0, \infty)$ and $\tau \in \mathbb{N}_0$, with $k_n = o(n)$, such that for all $n, m \in \mathbb{N}$,
 $a \in \mathcal{A}^n, b \in \mathcal{A}^m$

(UD) for all $\xi \in \mathcal{A}^{\tau}$

$$\mathbb{P}[a \xi b] \leq e^{k_n} \mathbb{P}[a] \mathbb{P}[b];$$

(SLD) there exists $0 \leq \ell(a, b) \leq \tau$ and $\xi(a, b) \in \mathcal{A}^{\ell}$ such that

$$\mathbb{P}[a \xi b] \geq e^{-k_n} \mathbb{P}[a] \mathbb{P}[b].$$

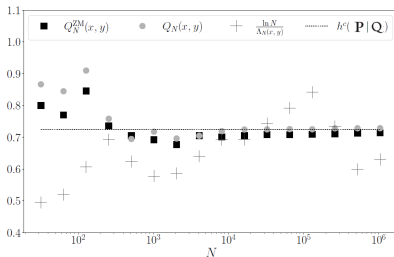
Step II: find new arguments which include a broader class of measures

\implies mZM: we replace
 $\Lambda_n(x, y) \mapsto \Lambda_n(x, y) + 1$
and

$$Q_n(x, y) := \frac{c_n(x, y) \log n}{n - c_n(x, y)}$$

Examples of decoupled measures:

- **Bernoulli measures;**
- **irreducible multi-level Markov chains;**
- equilibrium measures for 1D spin systems with summable interactions;
- g-measures;
- ψ -mixing measures;
- hidden-Markov measures with finite hidden alphabet;
- unravellings of irreducible repeated quantum measurements.



Theorems

Theorem (Cristadoro-Degli Esposti-Jakšić-Raquépas, 2023)

Let \mathbb{Q} satisfy (UD) and (SLD), and \mathbb{P} be ergodic. Then

$$\lim_{n \rightarrow \infty} \frac{\log W_n(x, y)}{n} = h^c(\mathbb{P}|\mathbb{Q}) \quad \text{for } (\mathbb{P} \otimes \mathbb{Q})\text{-almost every pair } (x, y) \text{ of sequences.}$$

Theorem (Barnfield-Grondin-P.-Raquépas, 2023)

Suppose that \mathbb{Q} satisfies (UD) and (SLD), and that \mathbb{P} is ergodic and satisfies (UD). Then, under mild decay conditions,

$$\lim_{n \rightarrow \infty} Q_n(x, y) = h^c(\mathbb{P}|\mathbb{Q}) \quad \text{for } (\mathbb{P} \otimes \mathbb{Q})\text{-almost every } (x, y) \in \mathcal{A}^{\mathbb{N}} \times \mathcal{A}^{\mathbb{N}}.$$

Recall that $Q_n(x, y)$ is analogous to $Q_n^{ZM}(x, y)$, but considering the *shortest* new word in x_1^n **not** found in y_1^n (as in the LZ77 algorithm!)

Some references

In the literature:

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Our results: Barnfield, Grondin, P. and Raquépas

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